### CHAPTER 5: PAIRWISE SEQUENCE ALIGNMENT AND DATABASE SEARCHING

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### Alignment algorithms (preview)

- Needleman-Wunsch (1970) and variations:
  - for aligning two sequences
  - uses dynamic programming to "consider" all possible alignments (10<sup>600</sup> for two sequences of length 1000!)
- FASTA: uses a heuristic method for efficient searches (though not guaranteed to find the optimal solution)
  - Creates dictionary of k-tuples for the query sequence which is checked against sequences in the database
  - A local alignment algorithm is used to complete the alignment
- BLAST (Basic Local Alignment Search Tool): also fast and uses a heuristic
  - Finds short matches (which do not have to be perfect)
  - Then uses local alignment to complete the alignment

#### Needleman and Wunsch Dynamic Programming method

- Dynamic programming method
  - A problem is broken down into smaller sub-problems that can be repeatedly (recursively) solved
- Consider the following alignment of the sequences

**x** = THISLINE and **y** = ISALIGNED

### **Dynamic Programming Concept**

The alignment below is the optimal global alignment, with BLOSUM-62 matrix, and linear gap penalty of 4:

A key observation is that the total alignment score is the sum of partial alignment scores:

Score:	-4	-4	+4	+4	-4	+4	+4	-4	+6	+5	-4	
	т	Н	I	S	-	L	I	– G	N	Е	_	
	-	_	I	S	A	L	I	G	N	E	D	
Score:	0					0		7				

For any part of the alignment: The score of the global alignment from positions  $1 \rightarrow L$  is the sum of the partial alignment scores from positions  $1 \rightarrow u$ ,  $u+1 \rightarrow v$ , and  $v+1 \rightarrow L$ 

#### Dynamic programming, key idea

The optimal alignment between sequences **x** and **y** can end in 3 possible ways (example for alignment between THIS and IS is shown):

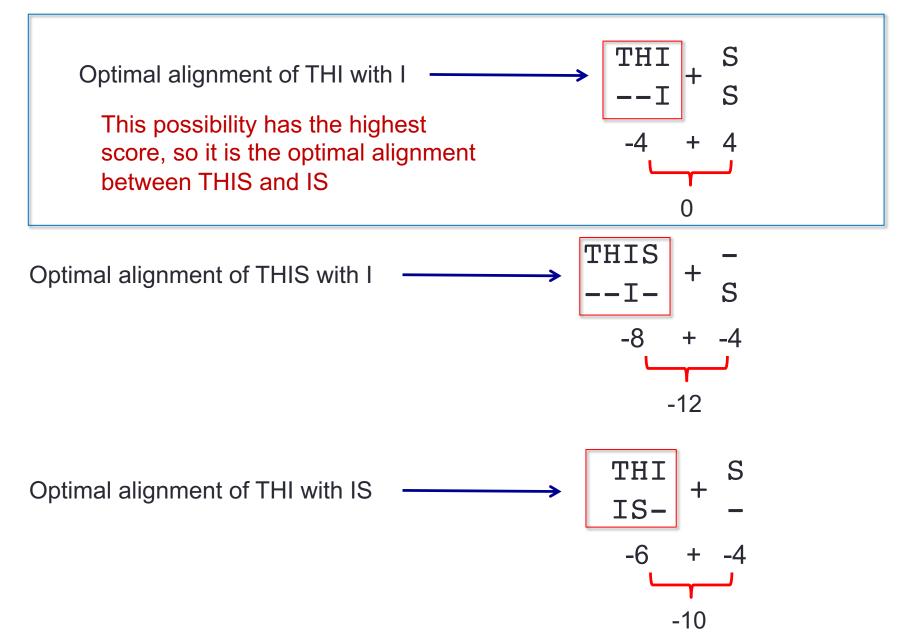
1. An alignment between the last character of  ${\boldsymbol x}$  and last character of  ${\boldsymbol y}$ 

2. A gap at the end of x aligned to the last character of y

3. A gap at the end of y aligned to the last character of x

S

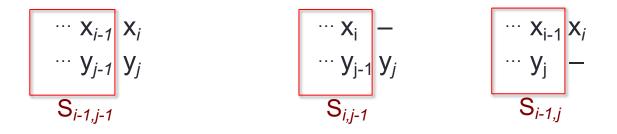
#### Three possibilities for optimal alignment ending btwn THIS and IS



#### Dynamic Programming Algorithm

Let  $S_{i,j}$  be the score of the optimal alignment of all characters up to  $x_i$  of sequence **x** and  $y_j$  of sequence **y**.

Then there are three possibilities for the end of this alignment



The optimal alignment score  $S_{i,j}$  is the maximum of the following

 $S_{i-1, j-1} + s(x_i, y_j)$   $S_{i, j-1} + g$   $S_{i-1, j} + g$ 

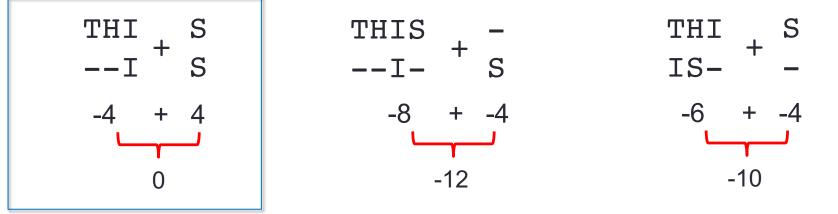
where  $s(x_i, y_j)$  is the substitution (or match/mismatch score) for character  $x_i$  aligned with  $y_j$ , and g is a constant gap penalty

Example:

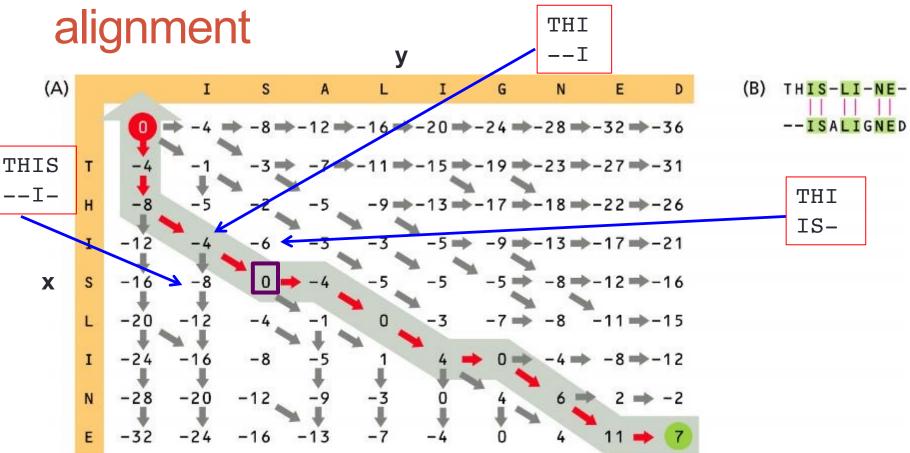
x = THISLINE and y = ISALIGNEDTHIS-LI-NE---ISALIGNED

Assume a constant (linear) gap penalty g = 4, and  $s(x_i, y_j)$  calculated using the BLOSUM-62 scoring matrix.

 $S_{4,2}$  is the optimal score for aligning THIS with IS and is the maximum of



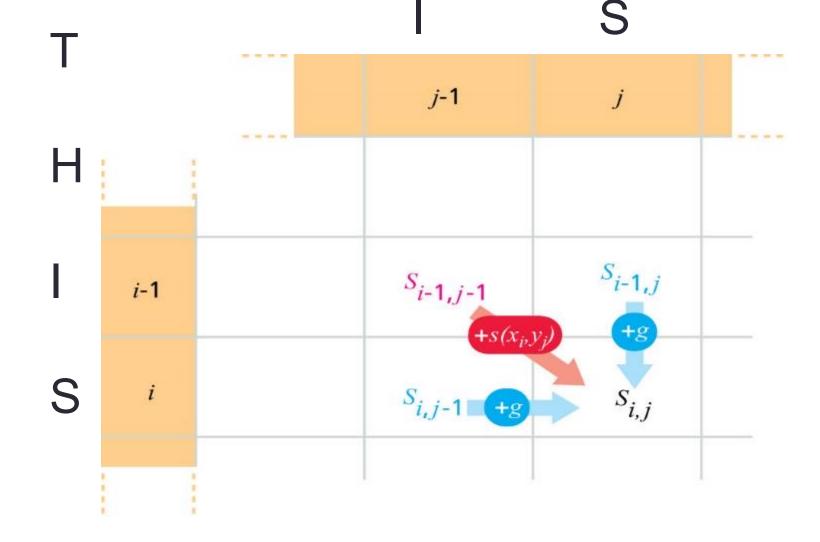
### Dynamic programming matrix for global



 $S_{i,j}$  is the score of the optimal alignment of all characters up to  $x_i$  of sequence **x** and  $y_i$  of sequence **y**.

The score of the optimal alignment of THIS with IS is  $S_{4,2} = 0$ .

Dynamic programming matrix illustrating possible optimal alignments for subsequences  $x_1, \ldots x_i$  and  $y_1, \ldots y_j$ 



# Algorithm for finding the optimal global alignment

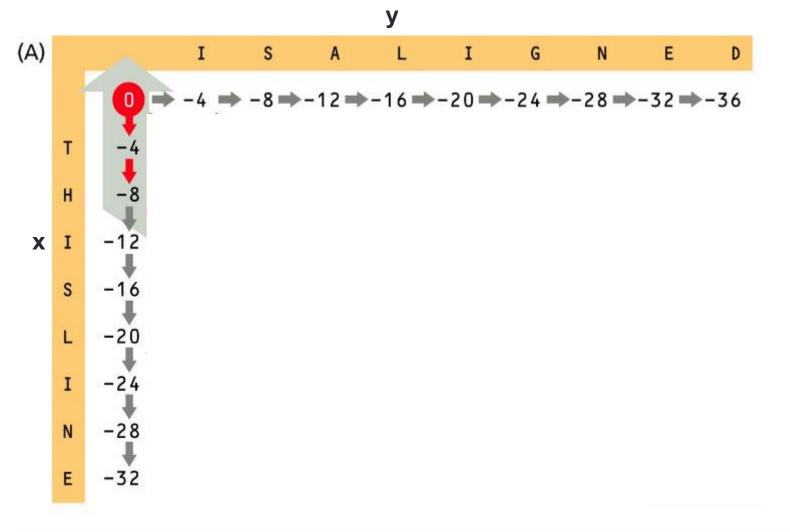
- For two sequences,  $x_1, \ldots x_m$  and  $y_1, \ldots y_n$
- Construct an (*m*+1) by (*n*+1) matrix (starting with element S<sub>0,0</sub> at the top left)
- For each *i*, the element  $S_{i,0}$  corresponds to the alignment

For each *j*, the element  $S_{0,j}$  corresponds to the alignment

- -  $\cdots$   $y_1$   $y_2$   $\cdots$   $y_i$ 

Fill in the matrix for all  $S_{i,0}$  and  $S_{0,i}$  appropriately

# Assignment of $S_{i,0}$ and $S_{0,j}$ when a linear gap penalty of 4 is used

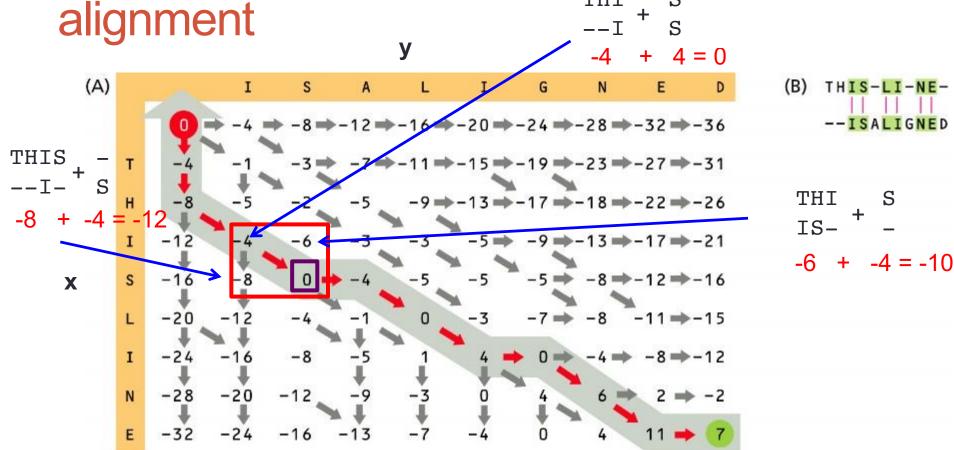


Algorithm for finding the optimal global alignment (con't)

set 
$$S_{i,j} = \max - \begin{cases} S_{i-1,j-1} + s(x_i, y_j) \\ S_{i-1,j} + g \\ S_{i,j-1} + g \end{cases}$$

and draw an arrow to  $S_{i,j}$  from the cell where the score was derived ( $S_{i-1,j-1}$ ,  $S_{i-1,j}$ , or  $S_{i,j-1}$ )

### Dynamic programming matrix for global

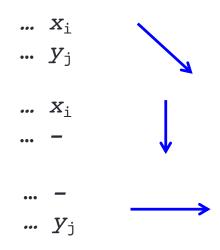


 $S_{i,j}$  is the score of the optimal alignment of all characters up to  $x_i$  of sequence **x** and  $y_i$  of sequence **y**.

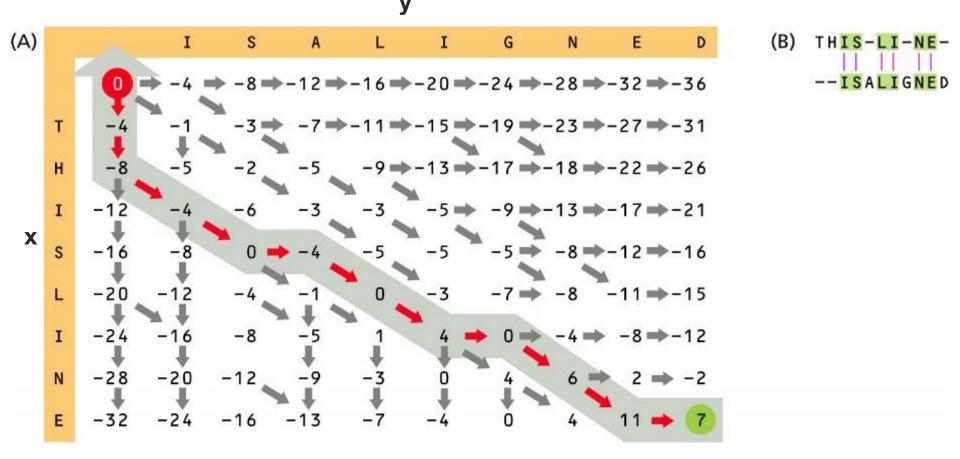
The score of the optimal alignment of THIS with IS is  $S_{4,2} = 0$ .

# Algorithm for finding the optimal global alignment (con't)

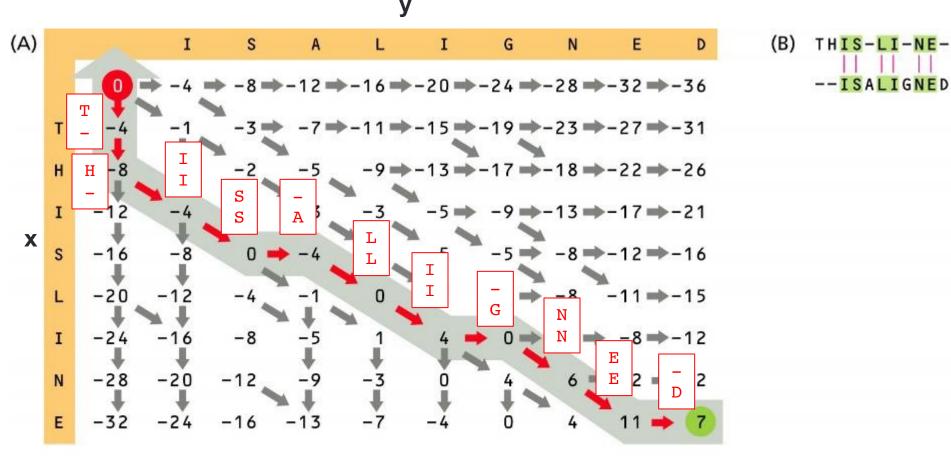
- After the dynamic programming matrix is filled in,  $S_{m,n}$  contains the score of the optimal global alignment
- A process called *traceback* is then used to recover the alignment itself, based on the path used to generate the optimal alignment score (based on the arrows), starting with  $S_{m,n}$  and working backwards to  $S_{0,0}$
- If the score  $S_{i,j}$  was derived
  - diagonally from  $S_{i-1,j-1}$ , then align  $x_i$  with  $y_j$
  - vertically from  $S_{i-1,j}$ , then align a gap with  $x_j$
  - horizontally from  $S_{i,j-1}$ , then align a gap with  $y_i$



# **Global alignment:** Dynamic programming matrix using linear gap penalty of 4



# **Global alignment:** Dynamic programming matrix using linear gap penalty of 4



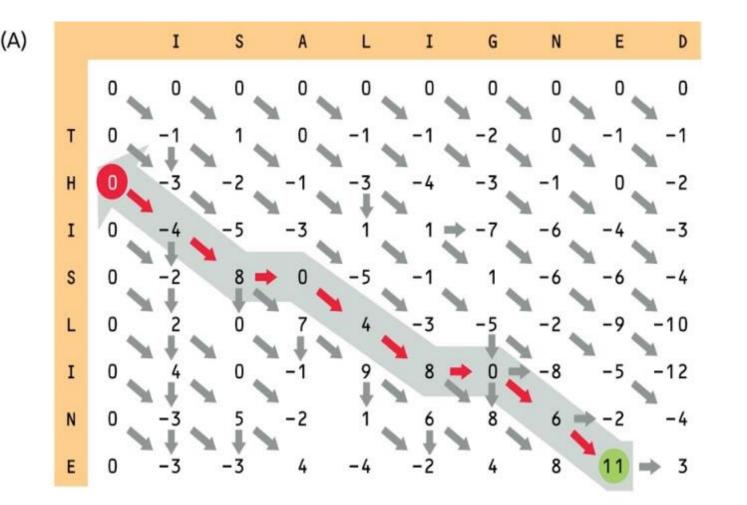
# **Global alignment:** Dynamic Programming Matrix using linear gap penalty of 8

(A)			I	S	А	L	I	G	N	E	D	(B)	THISLI <mark>NE</mark> -
		0	-8-	-16 🔿	-24 🔿	-32 🛋	-40 -	-48 🔿	-56 🔿	-64 🌩	-72		ISALIG <mark>NE</mark> D
	т	-8	-1	-7⇒	-15 🔿 ·	-23 🛋	-31	-39 🔿	-47	-55 🔿	-63		
	н	-16	-9	-2	-9	-17 🔿	-25	-33	-40	-47 🔿	-55		
	I	-24	-12	-10	-3	-7	-13 -	-21 🔿	-29 🔿	-37 🍑	-45		
	s	-32	-20	-8	-9	-5	-9	-13	-20	-28 🔿	-35		
	L	-40	-28	-16	-9	-5	-3	-11	-16	-23 🔿	-31		
	I	-48	-36	-24	-17	-7	-1	-7	-14	-19	-26		
	N	-56	-44	-32	-25	-15	-9	-1	-1 🔿	-9 ->	-17		
	E	-64	-52	-40	-33	-23	-17	-9	-1	4 🔿	-4		

# Modification of algorithm for semiglobal alignments

- A semiglobal alignment does not penalize for gaps at the beginning and end of an alignment
- Appropriate for global alignments of sequences that are not the same length
- Modifications to algorithm for semiglobal alignment:
  - Set  $S_{i,0}$  and  $S_{0,j}$  to 0 for all *i* and *j*
  - Start the traceback from the highest scoring element in the bottom row or last column

# **Semiglobal alignment:** Dynamic Programming Matrix using a linear gap penalty of 8



(B) THIS-LI-NE-|| || || --ISALIGNED

#### Modification of algorithm for local alignments (1<sup>st</sup> proposed by Smith-Watterman)

 A *local alignment* aligns *regions* of two sequences, and will not necessarily span the length of each sequence

 $X_i$ 

- Appropriate for identifying functional domains of a protein
- Modifications to algorithm for local alignment:

• set 
$$S_{i,j} = \max \begin{bmatrix} S_{i-1,j-1} + s(x_i, \\ S_{i-1,j} + g \\ S_{i,j-1} + g \\ 0 \end{bmatrix}$$

- Assumes that the expected alignment score is negative for random sequences and is positive for similar sequences
- Traceback starts from highest scoring matrix element anywhere in the matrix and ends at 0

# **Local alignment:** Dynamic Programming Matrix using a linear gap penalty of 8

